Problem Set 3
Let’s Simplex

AM121/ES121 — Fall 2017

Due 5:00 PM, Friday, September 29, 2017

Announcements

- The assignment is due by 5:00 PM, Friday, September 29, 2017.

Goals

This assignment has two goals. First, you will get practice solving linear programs using the Simplex method. Second, you will get a better understanding of linear programming through a better understanding of the Simplex method.

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But first, a distraction

In the last assignment, you got a bit of practice with modeling different objectives and constraints. But how do you know whether something is modelable via linear programming?

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Task 1

Consider the following constraint:

\[ \max(y_1, y_2) \geq 3 \]

In general, can we model this constraint with linear programming, without further knowledge of other constraints? (Hint: think about the geometry of the feasible region.)

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1 Making it rain

![Diagram of sprinkler system](image1)

Figure 1: Plant and sprinkler system diagrams. \( s \) indicates a sprinkler, \( p \) indicates a plant, and \( e \) indicates a piece of land without any plants.

![Spray delivery matrix](image2)

Figure 2: Spray delivered by an individual sprinkler

Taste Matters is growing impatient with all of your pen-and-paper theorizing—they want to see your work in action. They’ve thought about each of the formulations that you came up with in the previous assignment and have decided to keep it simple. They want you to minimize the amount of solution used for spraying the plants, subject to the same arrangement of sprinklers and plants as before (Figure 1). They also remind you that each sprinkler delivers a uniform spray covering a 3 by 3 region centered over the plant (Figure 2), and that each plant has a different minimum amount of solution required (specified in Table 1).

![Table of solution requirements](image3)

Table 1: Minimum amounts of solution required

For the following exercises, use the basic sprinkler model that you came up with in Task 7.1 of assignment 2. If you don’t remember it, feel free to use the formulation given in the solutions to assignment 2.
1. Implement your model in AMPL using the data provided by Taste Matters. Run the model to find the sprinkling policy that minimizes the amount of solution used to spray the plants.

2. Taste Matters is happy with the result that you’ve provided, but their engineering department has made a mistake: the actual output from each sprinkler is spread over the 3 by 3 region as follows:

<table>
<thead>
<tr>
<th></th>
<th>0.1</th>
<th>0.07</th>
<th>0.2</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.04</td>
<td>0.1</td>
<td>0.05</td>
<td></td>
</tr>
<tr>
<td>0.1</td>
<td>0.08</td>
<td>0.1</td>
<td></td>
</tr>
</tbody>
</table>

Taste Matters doesn’t think this mistake would nullify the policy you have found, but wants you to resolve the model at these rates to ensure that this is the case. Is it? (Hint: You should only have to change your data file).

End Task 2

2 Simplex by hand

Consider the following linear program:

\[
\begin{align*}
    \text{maximize} & \quad x_1 + x_2 + 2x_3 \\
    \text{subject to} & \quad x_1 \leq 3 \\
                         & \quad x_2 \leq 5 \\
                         & \quad x_1 + x_2 \leq 7 \\
                         & \quad -2x_1 + x_3 \leq 2 \\
                         & \quad x_1, x_2, x_3 \geq 0
\end{align*}
\]

The feasible region for this LP is shown below, with all extreme points labeled:

![Figure 3: Feasible region](image)
Task 3

1. Transform the linear program to standard equality form by introducing slack variables.

2. Solve the problem using the Simplex method, using the introduced slack variables as the initial feasible basis. In selecting an entering variable, choose the one with the smallest reduced cost. Record each tableau, feasible solution, and objective value you iterate through. As you solve the problem, indicate the path the Simplex method takes on the feasible region by referencing the labels on the diagram above. (Note: You could try doing pivots using MATLAB or Mathematica).

3. Is your solution a unique optimal solution? Why or why not?

End Task 3

Task 4

1. Consider changing the objective function for the linear program in the previous task from $x_1 + x_2 + 2x_3$ to $x_1 + 5x_2 + 2x_3$. Again solve the problem using the Simplex method, this time choosing the entering index based on the smallest subscript rule. Record each tableau, feasible solution, and objective value you iterate through. Be sure to reference the path the Simplex method takes about the feasible region by using the same labeled graph as in Task 3.

2. Is your solution a unique optimal solution? Why or why not?

End Task 4

3 Simplex, the works

Considering the following linear program:

$$\text{maximize} \quad -2x_1 + 4x_2 - x_3$$
$$\text{subject to} \quad 2x_1 + 3x_2 - 2x_3 \geq 4$$
$$5x_1 + 2x_2 - x_3 \leq 12$$
$$x_1, x_2 \geq 0; x_3 \text{ free}$$

Task 5

For this exercise, feel free to make use of MATLAB to perform the pivot operations as you have done on the previous assignment. This will help to avoid manual error!

1. Transform the linear program to standard equality form. This will require a number of steps.

2. Perform Phase I of the Simplex method to arrive at a feasible tableau. In selecting an entering variable, choose the one with the smallest subscript among the $x_i$’s before any other introduced variables. Be sure to record all of your steps.

3. Using the found tableau, solve the problem using Phase II of the Simplex method. In selecting an entering variable, choose the one with the smallest reduced cost. Record each tableau, feasible solution, and objective value you iterate through. You can error check by solving the problem in AMPL.

End Task 5
4 Learning from the tableau (by helping a friend...)

A friend of yours decides to play around with the following system of equations:

\[\begin{align*}
z & + x_1 + 2x_2 + 3x_3 + 2x_4 + x_5 + x_6 + x_7 = 0 \\
4x_1 + 2x_2 - 3x_3 + 5x_4 + x_5 & = 5 \\
5x_1 + 6x_2 - 2x_3 + 3x_4 & = 8
\end{align*}\]

There are also non-negativity constraints on all decision variables. He then performs a few pivots using MATLAB:

\begin{verbatim}
A = [1 0 0 0 0 1 1 0; 0 -1 2 3 2 1 0 0 5; 0 4 2 -3 5 0 1 0 8; 0 5 6 -2 3 0 0 1 4];
A = pivot(A, 4, 4);
A = pivot(A, 2, 5);
A = pivot(A, 3, 7);
\end{verbatim}

(Note: pivot(A, i, j) command needs to be defined in MATLAB as detailed in the pivoting guide here: http://am121.seas.harvard.edu/site/wp-content/uploads/2011/03/pivotingmaplelabmathematica1.pdf)

It then performs pivots on A resulting in A[i, j] = 1 and A[k, j] = 0 where k\(\neq \)i)

He arrives at the following tableau (the letters at the top of columns indicate the corresponding variables associated with coefficients in each row, with “=” left implicit between the penultimate and final columns.):

\[
\begin{bmatrix}
z & x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & b \\
1 & 4 & \frac{102}{13} & 0 & 0 & \frac{14}{13} & 0 & \frac{34}{13} & -\frac{15}{13} \\
0 & 1 & \frac{22}{13} & 0 & 1 & \frac{2}{13} & 0 & \frac{3}{13} & \frac{22}{13} \\
0 & -4 & -\frac{102}{13} & 0 & 0 & -\frac{1}{13} & 1 & -\frac{21}{13} & \frac{15}{13} \\
0 & -1 & \frac{12}{13} & 1 & 0 & \frac{3}{13} & 0 & -\frac{2}{13} & \frac{7}{13}
\end{bmatrix}
\]

But now your friend is clueless. Help.
Task 6

1. What are the basic variables? What is the corresponding basic solution?

2. Is the solution optimal for maximizing the value of $z$?

3. What does the solution imply about the solution to the following system of equations?

   \[
   \begin{align*}
   -x_1 + 2x_2 + 3x_3 + 2x_4 &= 5 \\
   4x_1 + 2x_2 - 3x_3 + 5x_4 &= 8 \\
   5x_1 + 6x_2 - 2x_3 + 3x_4 &= 4
   \end{align*}
   \]

End Task 6

Your friend thinks you are awesome and decides to play around with another linear program:

maximize $2x_3$

subject to

- $x_2 + x_3 \leq 8$
- $-x_2 + x_3 \leq 0$
- $x_1 + x_3 \leq 8$
- $-x_1 + x_3 \leq 0$
- $x_1, x_2, x_3 \geq 0$

Your friend adds slack variables $x_4, x_5, x_6,$ and $x_7$ as the basis variables for an initial feasible tableau (this now also labels the rows to indicate the “objective row” followed by the “constraints” with $x_4, x_5, x_6$ and $x_7$ isolated, respectively):

\[
\begin{bmatrix}
  z & x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & b \\
  z & 1 & 0 & 0 & -2 & 0 & 0 & 0 & 0 \\
  x_4 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\
  x_5 & 0 & 0 & -1 & 1 & 0 & 1 & 0 & 0 \\
  x_6 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 \\
  x_7 & 0 & -1 & 0 & 1 & 0 & 0 & 0 & 1 \\
\end{bmatrix}
\]

The basic solution is $\{0, 0, 0, 8, 0, 8, 0\}$, $z = 0$.

Your friend pivots by letting $x_3$ enter, leaving on $x_5$:

\[
\begin{bmatrix}
  z & x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & b \\
  z & 1 & 0 & -2 & 0 & 0 & 2 & 0 & 0 \\
  x_4 & 0 & 0 & 2 & 0 & 1 & -1 & 0 & 0 \\
  x_3 & 0 & 0 & -1 & 1 & 0 & 1 & 0 & 0 \\
  x_6 & 0 & 1 & 1 & 0 & 0 & -1 & 1 & 0 \\
  x_7 & 0 & -1 & 1 & 0 & 0 & -1 & 0 & 1 \\
\end{bmatrix}
\]

The basic solution is $\{0, 0, 0, 8, 0, 8, 0\}$, $z = 0$. “But this is exactly the same solution as I started with!” your friend says. “Help!”

You help by continuing the Simplex method, following the smallest subscript rule, yielding the following tableaux:
Task 7

1. Indicate the basic solution and objective values of the final two tableaux above. In addition, explain how the entering and leaving variables were chosen.

2. Sketch the feasible region in 3D. Trace the path that the Simplex algorithm took. Explain in a sentence why the same basic solution would show up more than once in this problem (Hint: think about the Simplex algorithm geometrically).

End Task 7

Having been so helpful to your friend, you decide to go the extra mile and come up with an example tableau to really help your friend understand the Simplex method. You have declared all variables as non-negative, and have left some elements of the tableau as lettered parameters (there is no normal meaning to any of these letter variables, i.e. don’t read “b” as the right-hand side or “z” as the objective.):

\[
\begin{bmatrix}
  z & x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & b \\
  z & 1 & -2 & 0 & 0 & 0 & 0 & 0 & 2 & 0 \\
x_4 & 0 & 2 & 0 & 0 & 1 & 1 & 0 & -2 & 8 \\
x_5 & 0 & -1 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\
x_6 & 0 & 2 & 0 & 0 & 0 & 0 & 1 & -1 & 8 \\
x_2 & 0 & -1 & 1 & 0 & 0 & -1 & 0 & 1 & 0 \\
  z & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 8 \\
x_1 & 0 & 1 & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & -1 & 4 \\
x_3 & 0 & 0 & 0 & 1 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 4 \\
x_6 & 0 & 0 & 0 & 0 & -1 & -1 & 1 & 1 & 0 \\
x_2 & 0 & 0 & 1 & 0 & \frac{1}{2} & -\frac{1}{2} & 0 & 0 & 4 \\
\end{bmatrix}
\]

You wish to explain to your friend the conditions under which the different types of scenarios occur in solving maximization linear programs. To avoid tarnishing your pristine reputation (to your friend, anyways), you review the conditions first yourself.

Task 8

For each of the following tasks, specify the conditions on the parameters that make the statement true. Aim to be as general as possible (that is, providing an example scenario is insufficient).

1. The current basic solution is optimal.
2. The current basic solution is infeasible.
3. The current basic solution is a degenerate basic feasible solution.
4. The current basic solution is feasible, but the objective value is unbounded.
5. The current basic solution is feasible, but can be improved by \( x_4 \) entering, and \( x_6 \) leaving the basis.
5  Termination

Without the issue of degeneracy, we can show that the Simplex algorithm terminates by using the guaranteed increase in the objective value at each iteration to show that no basis (from a finite set of bases) will be repeated. With degeneracy, we are not guaranteed that the objective value will increase at each iteration. Nevertheless, we can prove the following (and we do in the following exercise):

**Theorem 5.1.** Given an initial feasible tableau, the Simplex method using the smallest subscript rule will terminate in a finite number of iterations.

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**Task 9**

1. Place a bound on the maximum number of possible basic solutions one may encounter during the Simplex algorithm.

2. Given that two tableaux corresponding to the same basis can differ only in the order in which the equations are written, place a bound on the maximum number of possible tableaux one may encounter during the Simplex algorithm.

3. Deduce the theorem using the bound and Bland’s Theorem (which states that the Simplex method using the smallest subscript rule for entering and leaving variables cannot cycle).

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6  Turning in your assignment

You must turn in your assignment by 5:00 PM, Friday, September 29, 2017. For AMPL exercises, this includes any model and data files you have created. When writing down the solution from AMPL, always include both the objective value and the values assigned to variables (when the program is feasible and bounded, of course!)

Gather the AMPL model and data files you have created for this assignment, as well as a script file containing the AMPL commands you used to solve the problems. Please submit these files to the Canvax Dropbox at [https://canvas.harvard.edu/courses/13279/assignments](https://canvas.harvard.edu/courses/13279/assignments) before 5:00 PM, Friday, September 29, 2017. Please scan and submit your homework if you completed it on paper (make sure you submit the scanned portion as a single file). If you feel it may be hard to read, feel free to bring the original to class in addition to submitting online by the deadline.

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Congratulations on completing your third AM121 assignment!